

CRYPTOGRAPHY

Intro

Created by Alkalem, modified by Hanna3-14

```
import pwn

pwn.context.arch = "amd64"
pwn.context.os = "linux"

SHELLCODE = pwn.shellcraft.amd64.linux.echo('Test') + pwn.shellcraft
EXPLOIT = 0x45*b"\x90" + pwn.asm(SHELLCODE, arch="amd64", os="linux")

PROGRAM = b""
length = 20 + 16
for i in EXPLOIT:
    PROGRAM += i*b'+ ' + b'>'

    if i == 1:
        length += 5
    elif i > 1:
        length += 6
    length+= 13

(0x8000 - length) > 0x40:
    PROGRAM += b"<>"
    length += 2*13

    b"."["
    (0x8000 - length) + 7 -1
    (F+0x10)*b"<"

(host", 1337) as conn:
    (b"Brainf*ck code: ")
    PROGRAM)
    e()
```

OVERVIEW

- security goals
- cryptographic ciphers/protocols
 - classical cryptography
 - randomness
 - symmetric cryptography
 - asymmetric cryptography
- typical vulnerabilities
- tools
- further topics
- tasks

SECURITY GOALS

- **confidentiality:** protecting personal data
- **integrity:** data has not been modified
- **authenticity:** exactly what is claimed

CLASSICAL CRYPTOGRAPHY

CAESAR CIPHER

- each letter is shifted by a fixed value K
- $encrypt_K(P) = (P + K) \bmod 26$
- $decrypt_K(C) = (C - K) \bmod 26$

Example for $K = 23$:

Plaintext:	T	H	E	Q	U	I	C	K	B	R	O	W	N	F	O	X	...
Ciphertext:	Q	E	B	N	R	F	Z	H	Y	O	L	T	K	C	L	U	...

CLASSICAL CRYPTOGRAPHY

VIGENÈRE CIPHER

- choose a key of multiple letters, shift each letter according to the key letter
- attacks: determine key length, frequency analysis, (Kasiski examination)

Example:

Plaintext:	A	T	T	A	C	K	A	T	D	A	W	N
Key:	L	E	M	O	N	L	E	M	O	N	L	E
Ciphertext:	L	X	F	O	P	V	E	F	R	N	H	R

CLASSICAL CRYPTOGRAPHY

ONE TIME PAD (OTP)

- use key with same length as plaintext
- **information-theoretically secure**, if key is chosen equally distributed at random and used only once
- key exchange needs to be done separately via a secure channel
- mostly not realizable

RANDOMNESS

- in most programming languages default pseudo-random number generators (PRNGs) are not cryptographically secure
 - ⇒ state can be recovered
- cryptographically secure RNGs:
 - /dev/urandom
 - hardware RNGs
 - RNGs of the cryptographic libraries (e.g., secrets or Crypto.Random in python)

SYMMETRIC CRYPTOGRAPHY

STREAM CIPHERS

- pseudo-random key stream generated from key with an PRNG
- key stream is XORed with plaintext stream
- examples: RC4, SEAL, Salsa, CryptMT
- attacks:
 - known plaintext: calculate parts of the key stream when parts of plaintext are known
 - key reuse: two messages are encrypted with same key stream, difference (XOR) between plaintexts is observable

SYMMETRIC CRYPTOGRAPHY

BLOCK CIPHERS

- encrypt blocks of fixed length
- padding: extend messages to full block length
- examples: DES, IDEA, RC5, AES, Blowfish, ...
- modes of operation: (e.g. ECB, CBC, CTR, GCM)
- attacks:
 - against cipher: differential or linear cryptanalysis
 - different attacks against different modes of operation

RSA - KEY GENERATION

- choose large prime numbers p and q
- calculate the modulus $N = p * q$
- calculate $\phi(N) = (p - 1) * (q - 1)$
- choose e with $\gcd(e, \phi(N)) = 1 \wedge 1 < e < \phi(N)$
- calculate d as inverse of e under modulus $\phi(N)$

$$e * d \equiv 1 \pmod{\phi(N)}$$

- public key: N, e
- private key: d

RSA - ENCRYPTION AND DECRYPTION

- encryption: $c = m^e \pmod N$
- decryption: $c^d = (m^e)^d = m^{ed} \pmod{\phi(N)} = m^1 \pmod N$
- RSA without special padding is homomorphic
($Enc(m_1, pk) * Enc(m_2, pk) = Enc(m_1 * m_2, pk)$) and deterministic
- use RSA-OAEP if that is problematic

RSA - ATTACKS

- factoring N has complexity of about $\exp(\log(N)^{\frac{1}{3}} (\log\log N)^{\frac{2}{3}})$, infeasible for reasonable choice of N
- in some cases, attacks in polynomial time possible:
 - small private exponent d ($d < \frac{1}{3} N^{\frac{1}{4}}$): Wiener's attack
 - for small public exponent or partially known prime factor: Coppersmith's attack
 - $m < N^{\frac{1}{e}}$: calculate message as root of ciphertext
 - message sent to many recipients using same public exponent: Hastad's Broadcast Attack

ELLIPTIC CURVES

- elliptic curve equation:
- group: generator point G , point addition and multiplication with natural number
- cyclic EC group over $\mathbb{Z}_p, p > 3$
- point (x, y) on curve iff $y^2 \equiv x^3 + ax + b \pmod{p}$, plus (imaginary) point at infinity $O, a, b \in \mathbb{Z}_p$ with $4a^3 + 27b^2 \not\equiv 0 \pmod{p}$
- harder to attack, can use smaller keys for same security level

TYPICAL VULNERABILITIES

- implementation mistake: incorrect/vulnerable custom implementation, incorporated incorrectly into application
- conceptual mistake: incorrect use or not sufficient for use case
- theoretic mistake: violated condition for security, advanced maths or theoretic computer science necessary, "read the paper"
- well-known and documented attacks (e.g. length extension attack)
- oracles

TOOLS

- CyberChef
- <https://factordb.com/>
- [sagemath](#) (free open-source mathematics software system)
- [Z3](#) (theorem prover)

FURTHER TOPICS

- post-quantum cryptography
- pairing-based cryptography
- zero knowledge

START PLAYING CRYPTO CTF

- <https://intro.kitctf.de/>
- other platforms:
 - <https://cryptohack.org/> (Easy to hard, with good explanations)
 - <https://cryptopals.com/> (Implement cryptosystems and attacks)
 - <https://overthewire.org/wargames/krypton> (Classical crypto)
 - <https://imaginaryctf.org/> (not only crypto but also other CTF challenges)